

# **Non-Intel Tools at Claix**

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## **Allinea Performance Reports**





## Simple to use tool to get an performance overview mpiexec -> perf-report mpiexec ...



### Summary: jacobi.exe is Compute-bound in this configuration



Time spent running application code. High values are usually good. This is **high**; check the CPU performance section for advice

Time spent in MPI calls. High values are usually bad. This is **low**; this code may benefit from a higher process count

Time spent in filesystem I/O. High values are usually bad. This is **negligible**; there's no need to investigate I/O performance

This application run was Compute-bound. A breakdown of this time and advice for investigating further is in the CPU section below. As little time is spent in MPI calls, this code may also benefit from running at larger scales.

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# **Allinea Performance Reports**

### CPU

A breakdown of the 82.0% CPU time: Single-core code 2.9% | OpenMP regions 97.1% Scalar numeric ops 41.0% Vector numeric ops 0.0% | Memory accesses 59.0%

The per-core performance is memory-bound. Use a profiler to identify time-consuming loops and check their cache performance.

No time is spent in vectorized instructions. Check the compiler's vectorization advice to see why key loops could not be vectorized.

### 1/0

A breakdown of the 0.0% I/O time:

Time in reads	0.0%	6
Time in writes	0.0%	E
Effective process read rate	0.00 bytes/s	ß
Effective process write rate	0.00 bytes/s	I.

No time is spent in I/O operations. There's nothing to optimize here!

#### Memory

Per-process memory usage may also affect scaling:

Mean process memory usage	251 MiB	2	
Peak process memory usage	263 MiB		
Peak node memory usage	4.0%	1	

The peak node memory usage is very low. Running with fewer MPI processes and more data on each process may be more efficient.

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## MPI

A breakdown of the 18.0% MPI time:

Time in collective calls	69.6 <b>%</b>	
Time in point-to-point calls	30.4%	
Effective process collective rate	6.65 kB/s	1
Effective process point-to-point rate	83.3 MB/s	

Most of the time is spent in collective calls with a very low transfer rate. This suggests load imbalance is causing synchronization overhead; use an MPI profiler to investigate.

The point-to-point transfer rate is low. This can be caused by inefficient message sizes, such as many small messages, or by imbalanced workloads causing processes to wait.

### OpenMP

A breakdown of the 97.1% time in OpenMP regions:

Computation	86.9%	
Synchronization	13.1%	1
Physical core utilization	22.2%	1
System load	25.0 <mark>%</mark>	

Physical core utilization is low and some cores may be unused. Try increasing OMP\_NUM\_THREADS to improve performance.







## Score-P



- Supports Tracing and Profiling
- Uses direct instrumentation
- It also supports C/C++ and Fortran with MPI, OpenMP and hybrid codes.
- Useful for large scale applications.

## Usage:

- Precede your compiler command with scorep
   icc test.c –openmp –o a.out → scorep icc test.c –openmp –o a.out
   mpicc test.c –openmp –o a.out → scorep mpicc test.c –openmp –o a.out
- 2.a Run your application as usual to generate a profile
- 2.b Set SCOREP\_ENABLE\_TRACING=true, SCOREP\_ENABLE\_PROFILING=false and run the application for a <u>trace</u>.
- 3. Analyze the data in scorep-XXXXXX

## **Performance API (PAPI)**



## Standard API to access hardware counters

- → provides access to a set of hardware counters with standardized names and over a standardized interface
- $\rightarrow$  used in many tools for hardware counter access (also in Score-P)
- → papi\_avail provides a list of available counters
- → allows also to measure counters to get MFLOPS

Name	Code	Avail	Deriv	Description (Note)
PAPI_L1_DCM	0x80000000	Yes	No	Level 1 data cache misses
PAPI_L1_ICM	0x80000001	Yes	No	Level 1 instruction cache misses
PAPI L2 DCM	0x80000002	Yes	No	Level 2 data cache misses
PAPI L2 ICM	0x80000003	Yes	No	Level 2 instruction cache misses
PAPI L3 DCM	0x80000004	No	No	Level 3 data cache misses
PAPI L3 ICM	0x80000005	No	No	Level 3 instruction cache misses
PAPI L1 TCM	0x80000006	Yes	Yes	Level 1 cache misses
PAPI L2 TCM	0x80000007	Yes	No	Level 2 cache mis
				developed at:



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# **Vampir – Function Summary**



- Overview of the total time spend in functions.
- Time spend in MPI with sending or receiving messages can be seen.
   Change Event Category to "Function" to split the "Function Groups" and get more details.

	Function Summary										
All Processes, Accumulated Exclusive Time per Function Group											
11 s	10 s	9 s	8 s	7 s	бs	5 s	4 s	3 s	2 s	1 s	0 s
11.365 s											OMP-LOOP
									1.93	бs	OMP-SYNC
									1.7	764 s	MPI
										¢.301	s 🔤 OMP-PREG
										68.823	ms Application
										<10	) ms VT_API
										<	Lms OMP
1	1				1				i	i	1

## Vampir - Timeline

- The Timeline gives a detailed view of all events.
- Regions and Messages of all Processes and Threads are shown.
  - Zoom horizontal or vertical for more detailed information.
- Click on a message or region for specific details.



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# **Vampir – Communication Matrix**

- MPI Communication Matrix for Point to Point Messages
- Overview over the communication behavior of the application
- Coupled with timeline view

## Different Views:

- Aggregated message
   volume
- → Min/Max message size
- → Average message size
- Min/Max/Avg transfer time
- → ....



## Vampir – Process Timeline



The callstack can give even more information on the functions called on every thread/process.



## Scalasca – Cube Data Browser





- 1. Metric tree
- 2. Call tree

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3. Topology tree

All views are coupled from left to right: 1. choose a metric

-> this metric is shown for all functions

choose a function
 the right view
 shows the distribution
 over processes





## Scalasca – Call Tree





## Automatic Pattern Detection with Scalasca



- For large scale applications visualizing traces might be to much information.
- Aggregating everything in a profile might lose important information
- Scalasca allows to search for performance problems automatically.
- By rerunning the trace and comparing the time stamps, several situations can be automatically detected.

Example: The late-sender pattern



The time lost through this problem (red arrow) is accumulated over the complete run and stored in a profile.

## Automatic Pattern Detection with Scalasca



# Many patterns can be detected, like: Computational load imbalance:



# OpenMP management time:



## **Scalasca - Traces**



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