

Programming OpenMP

Cut-off strategies

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Example: Sudoku revisited

Parallel Brute-force Sudoku



This parallel algorithm finds all valid solutions

	_					_	<u>J</u>	_			_				
	6						8	11			15	14			16
15	11				16	14				12			6		
13		9	12					3	16	14		15	11	10	
2		16		11		15	10	1							
	15	11	10			16	2	13	8	9	12				
12	13			4	1	5	6	2	3					11	10
5		6	1	12		9		15	11	10	7	16			3
	2				10		11	6		5			13		9
10	7	15	11	16				12	13						6
9						1			2		16	10			11
1		4	6	9	13			7		11		3	16		
16	14			7		10	15	4	6	1				13	8
11	10		15				16	9	12	13			1	5	4
		12		1	4	6		16				11	10		
		5		8	12	13		10			11	2			14
3	16			10			7			6				12	

(1) Search an empty fie

(2) Try all numbers:

(2 a) Check Sudoku

If invalid: skip

If valid: Go to ne #pragma omp task field

first call contained in a #pragma omp parallel #pragma omp single such that one tasks starts the execution of the algorithm

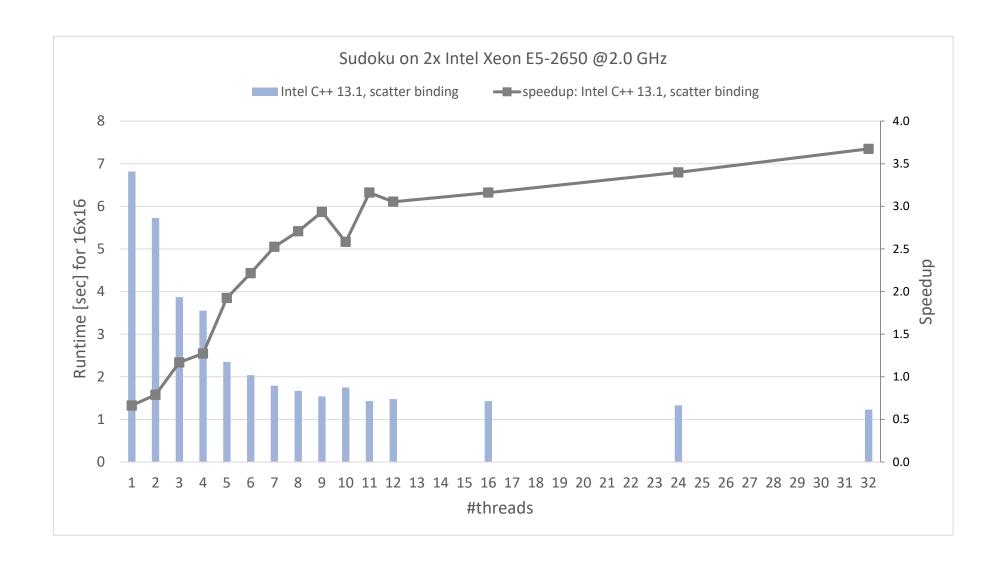
needs to work on a new copy of the Sudoku board

Wait for completion

#pragma omp taskwait wait for all child tasks

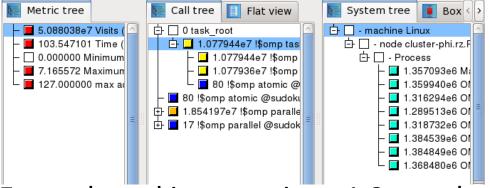
Performance Evaluation



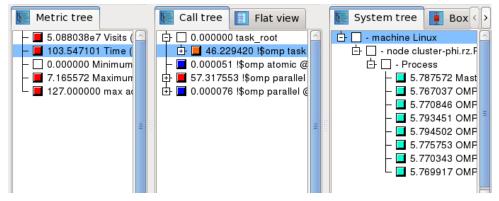


Performance Analysis

Event-based profiling provides a good overview :



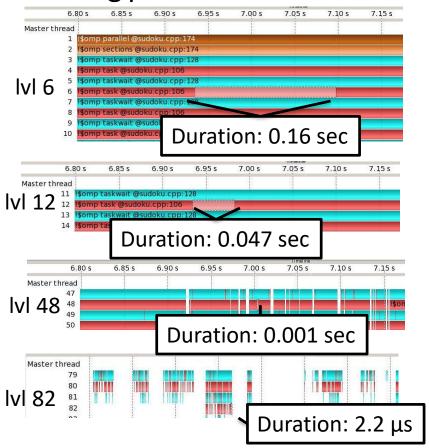
Every thread is executing ~1.3m tasks...



- ... in ~5.7 seconds.
- => average duration of a task is \sim 4.4 µs



Tracing provides more details:

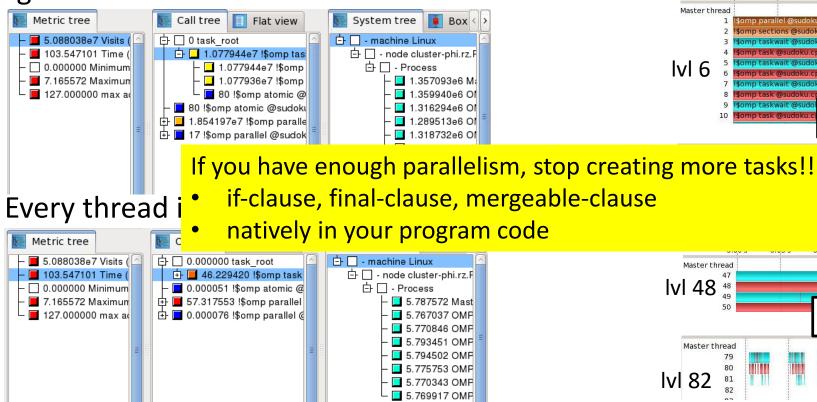


Tasks get much smaller down the call-stack.

Performance Analysis



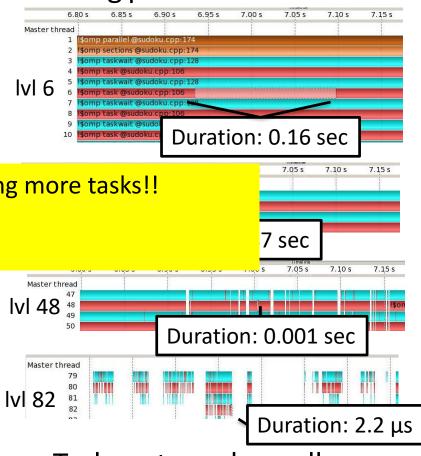
Event-based profiling provides a good overview :



... in ~5.7 seconds.

=> average duration of a task is \sim 4.4 µs

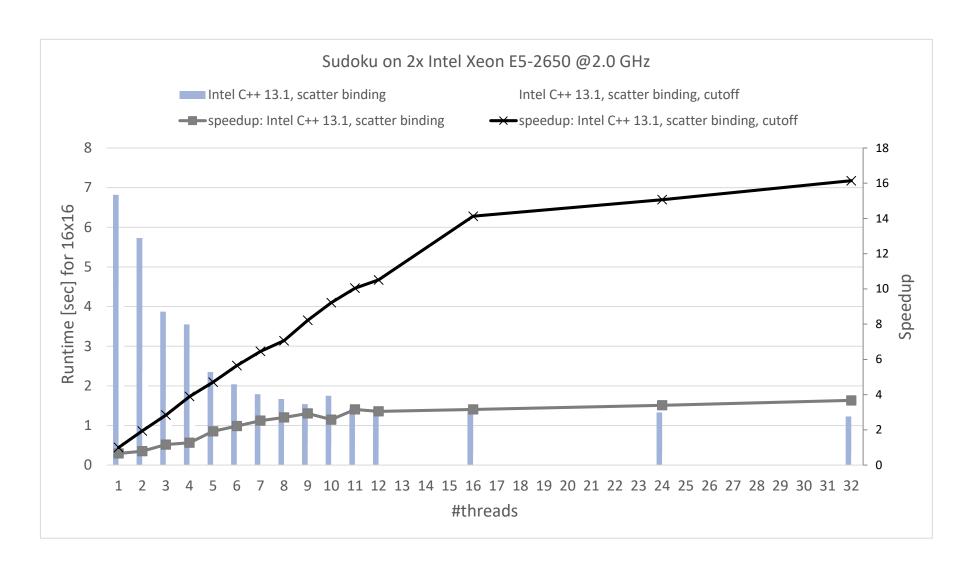
Tracing provides more details:



Tasks get much smaller down the call-stack.

Performance Evaluation (with cutoff)







Improving Tasking Performance: Cutoff clauses and strategies

The if clause



- Rule of thumb: the if (expression) clause as a "switch off" mechanism
 - → Allows lightweight implementations of task creation and execution but it reduces the parallelism
- If the expression of the if clause evaluates to false
 - → the encountering task is suspended
 - → the new task is executed immediately (task dependences are respected!!)
 - the encountering task resumes its execution once the new task is completed
 - → This is known as undeferred task

```
int foo(int x) {
  printf("entering foo function\n");
  int res = 0;
  #pragma omp task shared(res) if(false)
  {
     res += x;
  }
  printf("leaving foo function\n");
}
```

Really useful to debug tasking applications!

■ Even if the expression is false, data-sharing clauses are honored

The final clause

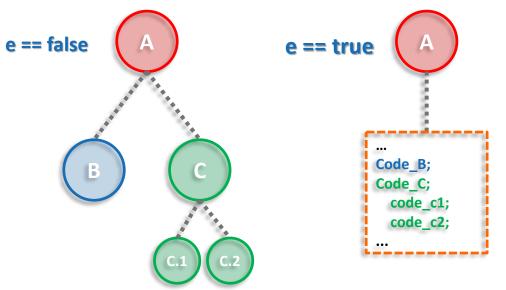


- The final (expression) clause
 - → Nested tasks / recursive applications
 - → allows to avoid future task creation → reduces overhead but also reduces parallelism
- If the expression of the final clause evaluates to true
 - → The new task is created and executed normally but in its context all tasks will be executed immediately

by the same thread (included tasks)

```
#pragma omp task final(e)
{
    #pragma omp task
    { ... }
    #pragma omp task
    { ... #C.1; #C.2 ... }
    #pragma omp taskwait
}
```

Data-sharing clauses are honored too!



The mergeable clause



- The mergeable clause
 - → Optimization: get rid of "data-sharing clauses are honored"
 - → This optimization can only be applied in undeferred or included tasks
- A Task that is annotated with the mergeable clause is called a mergeable task
 - → A task that may be a merged task if it is an undeferred task or an included task
- A merged task is:
 - → A task for which the data environment (inclusive of ICVs) may be the same as that of its generating task region
- A good implementation could execute a merged task without adding any OpenMPrelated overhead
 Unfortunately, there are no OpenMP

Unfortunately, there are no OpenMP commercial implementations taking advantage of final neither mergeable =(